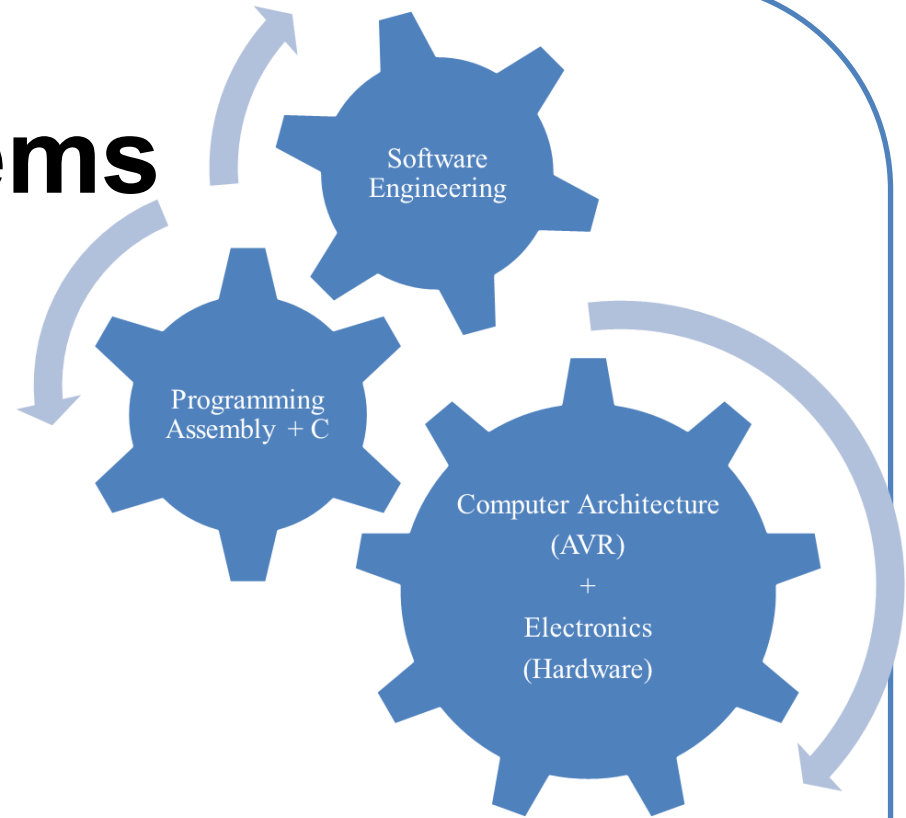


# Embedded Systems



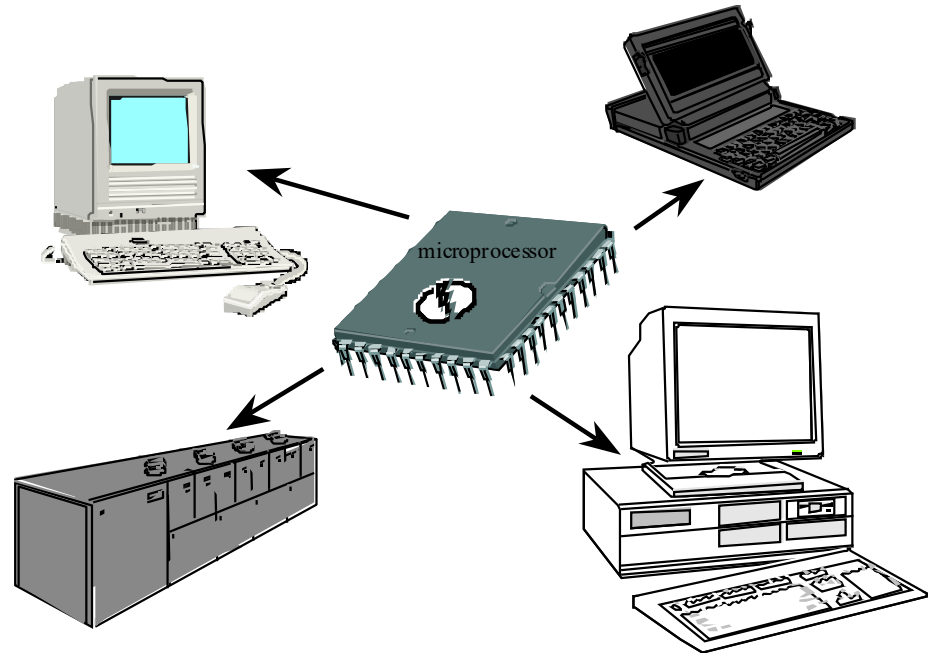
- Introduction to Embedded System.
- Review of Electronics Part (KCL, KVL, Parallel/Series Resistance, DAC and ADC).
- Review for Digital Systems ( Binary Number, Logic, MOS Implantation, Computer Architecture).
- Intro to Programming Concepts (Structure and Concurrent).

# What ? Embedded Systems

- **An embedded system is an **electronic system** that:**
  - includes a microcomputer embedded or **hidden** inside.
  - has software programmed into **ROM**.
  - has software that is **not accessible to the user** of the device
  - is configured to perform a **specific** dedicated application (software solves only a limited range of problems )

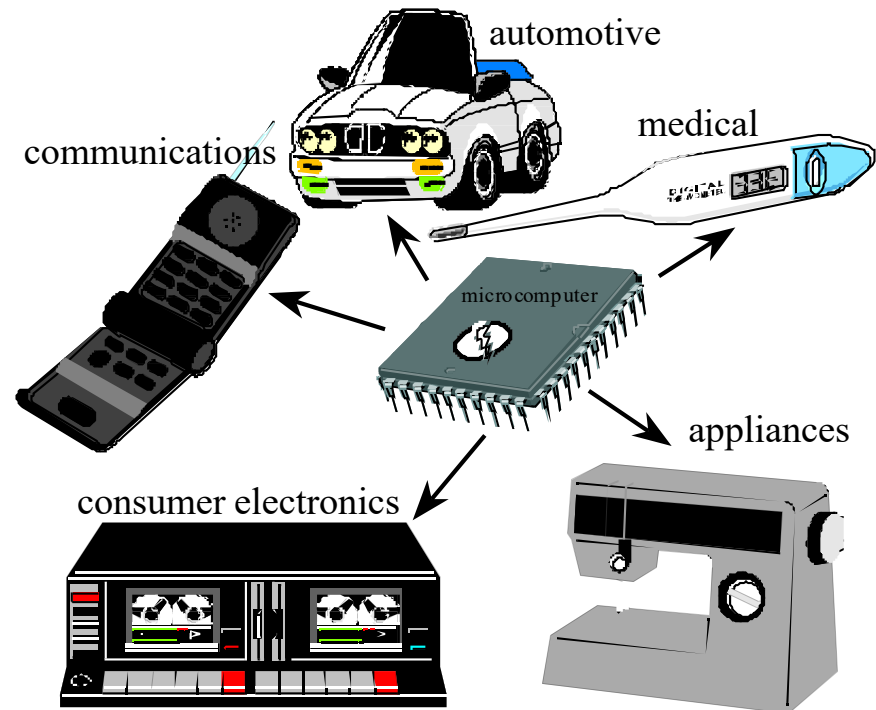
# General Purpose System

- keyboard
- disk
- graphics display
- software useful for a wide variety of purposes
- software that can be changed by user



# Embedded System

- Accepts inputs,
- Performs calculations
- Generates outputs
- Runs in “real time.”

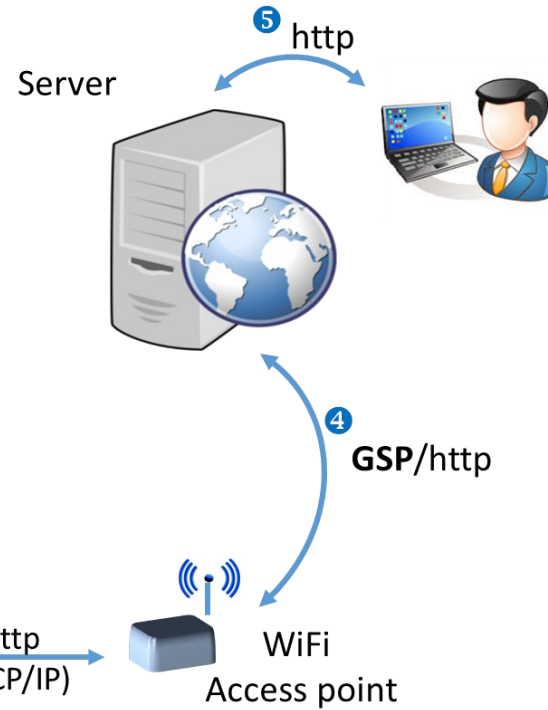
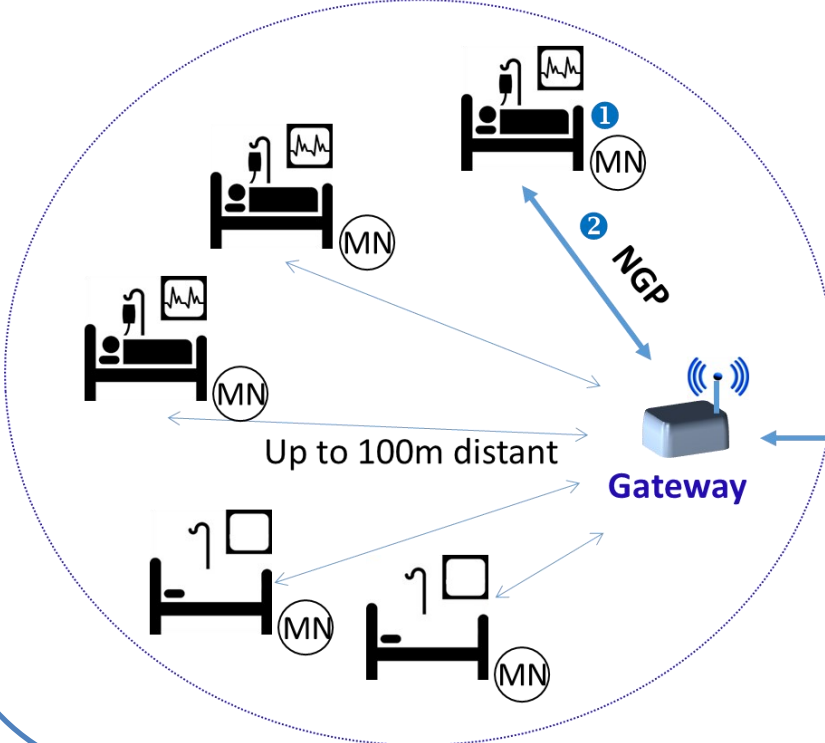


## The **internet of things** (IoT)

- Embedded systems (sensors and actuators).
- Internet.

# ICU Utilization System

Gateway coverage zone  
inside Intensive Care Room



NGP = Node Gateway Protocol  
GSP = Gateway Server Protocol

(MN) Monitoring Node – watch and report the activity of ICU

# Embedded Systems “Big Ideas”

- **HW/SW Architecture**
  - Non processor centric view of architecture
  - Microcontroller, FPGA, analog circuits
- **Bowels of the “operating system”**
  - Specifically, the lower half of the OS
  - Concurrency, parallelism, synchronization
- **Real world design**
  - performance vs. cost tradeoffs, constraints
- **Analyzability**
  - how do you “know” that your drive-by-wire system will function correctly?
- **Application-level techniques**
  - Power Aware Programming

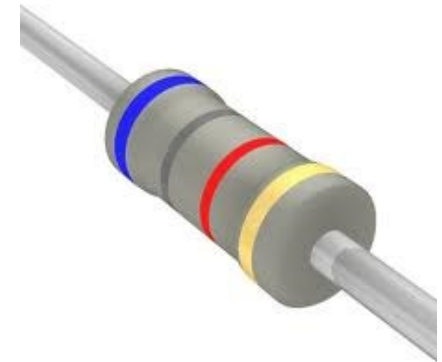
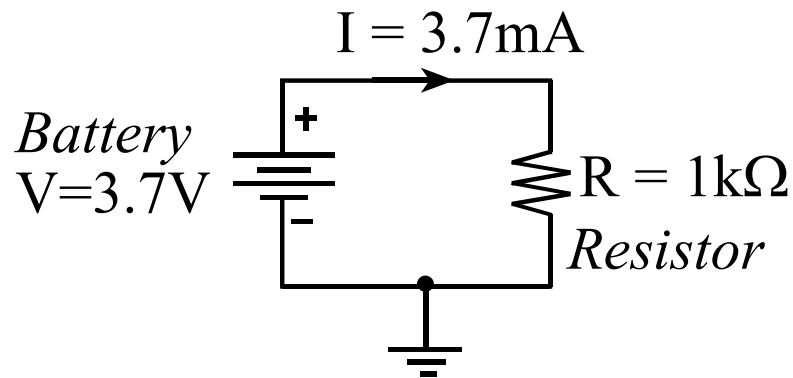
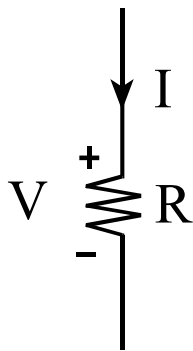
# Review of Electronics

## *Ohm's Law*

$$V = I * R \quad \text{Voltage} = \text{Current} * \text{Resistance}$$

$$I = V / R \quad \text{Current} = \text{Voltage} / \text{Resistance}$$

$$R = V / I \quad \text{Resistance} = \text{Voltage} / \text{Current}$$



$$P = V * I \quad \text{Power} = \text{Voltage} * \text{Current}$$

$$P = V^2 / R \quad \text{Power} = \text{Voltage}^2 / \text{Resistance}$$

$$P = I^2 * R \quad \text{Power} = \text{Current}^2 * \text{Resistance}$$

# Review of Electronics (Cont...)

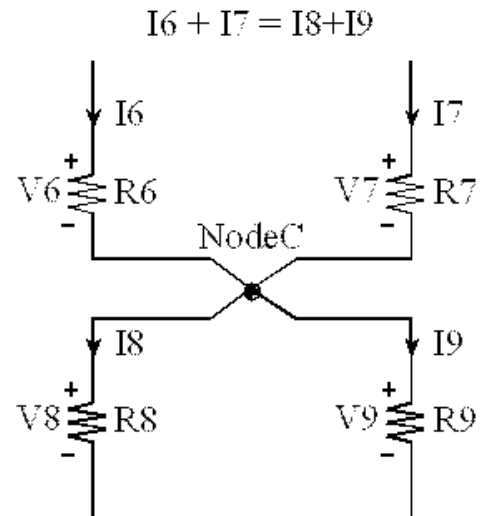
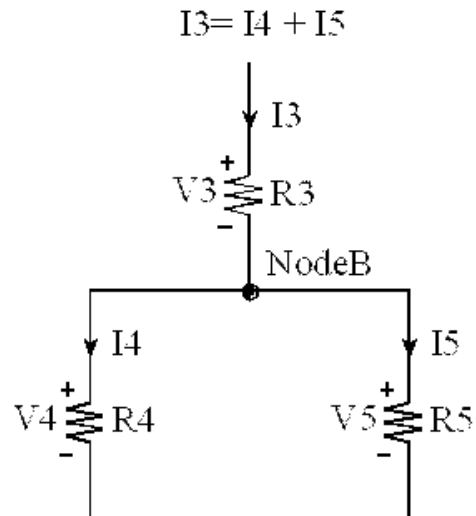
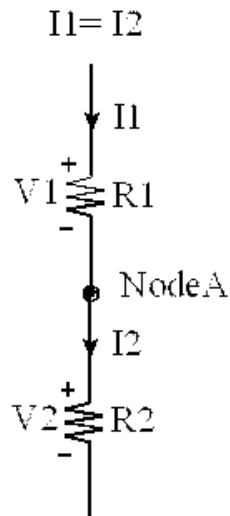
- **Voltage**: an electrical potential.
- **Current**: the flow of charge (electrons)
- **Power**: the rate of energy change.
- **Energy**: defines the amount of work that can be done
- **Resistance**: potential divided by flow



# Review of Electronics (Cont...)

## *Kirchhoff's Current Law (KCL).*

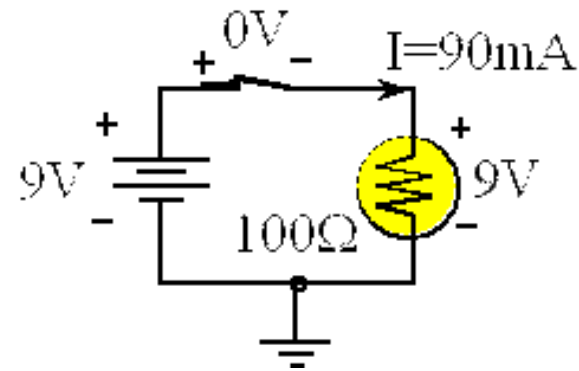
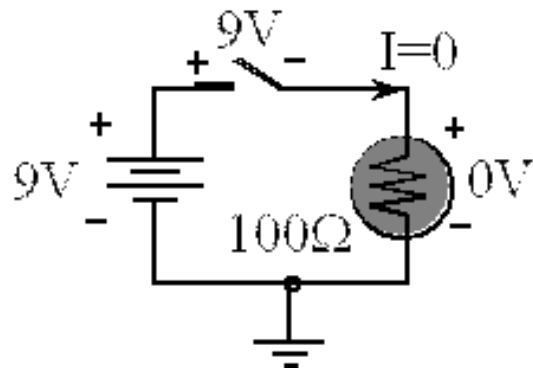
The sum of the currents **into a node** equal the sum of the currents **leaving the node**.



# Review of Electronics (Cont...)

*Kirchhoff's Voltage Law (KVL).*

The sum of the voltages **around the loop** is zero.



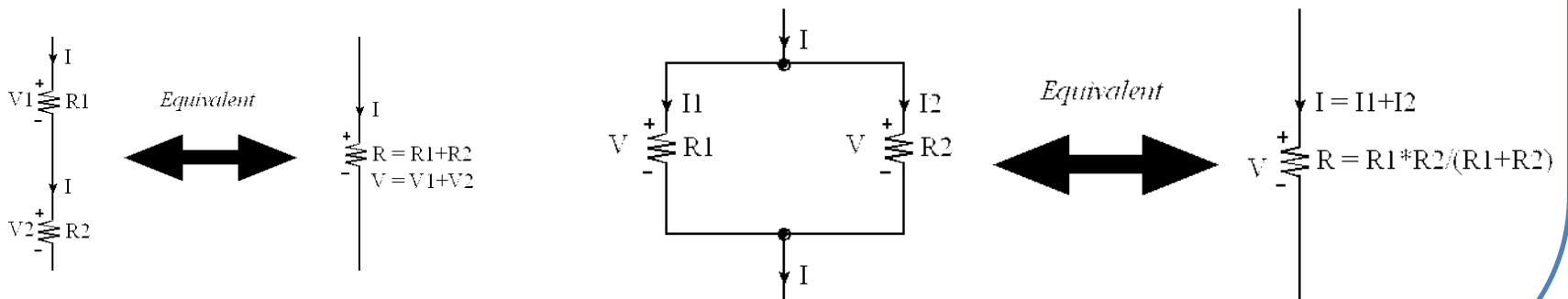
# Review of Electronics (Cont...)

## *Series resistance*

If resistor **R1** is in series with resistor **R2**, this combination behaves like one resistor with a value equal to **R1+R2**

## *Parallel resistance*

If resistor **R1** is in parallel with resistor **R2**, this combination behaves like one resistor with a value equal to **R1\*R2/(R1+R2)**



# Review of Electronics (Cont...)

Consider this 3-bit digital to analog converter.

Define a 3-bit number  $n$  (0 to 7) which specifies the three switch positions.

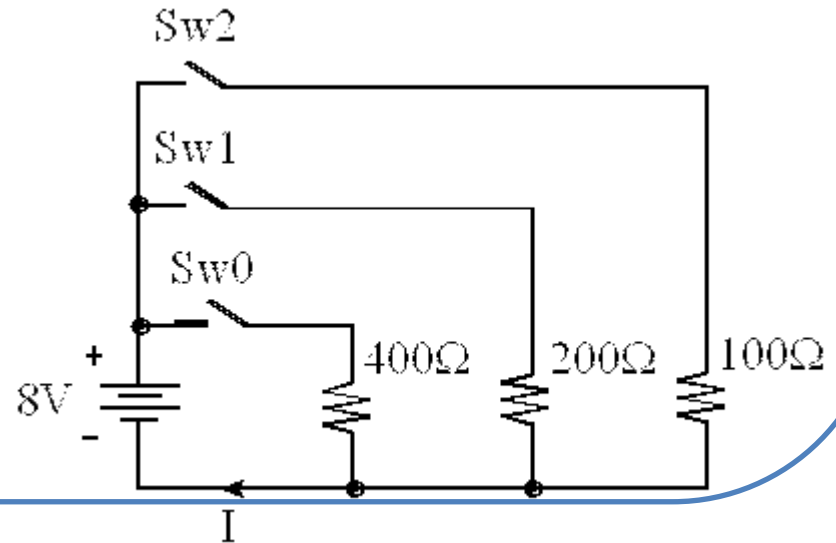
$n = 0$  means none are pushed.  $n = 1$  means Sw0 is pushed.

$n = 2$  means Sw1 is pushed.  $n = 3$  means Sw1 and Sw0 are pushed.

$n = 4$  means Sw2 is pushed.  $n = 5$  means Sw2 and Sw0 are pushed.

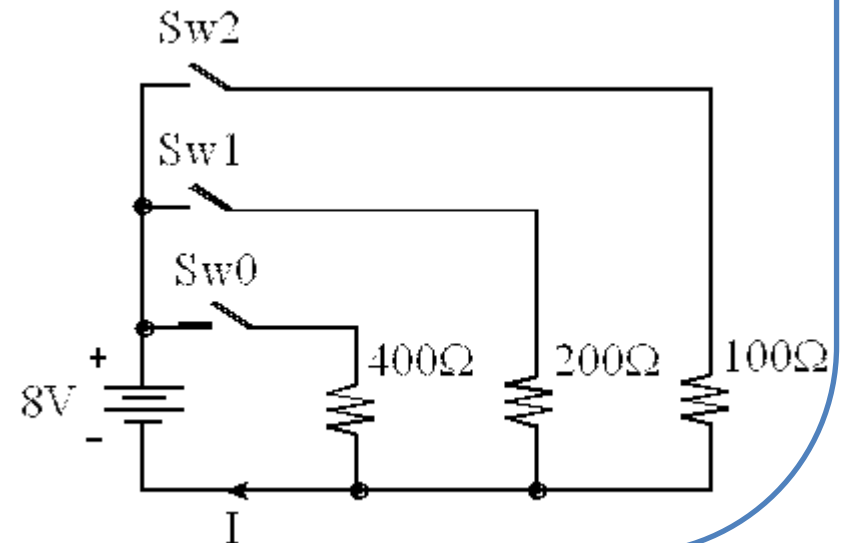
$n = 6$  means Sw2 and Sw1 are pushed.  $n = 7$  means all are pushed.

Derive a relationship between the current  $I$  and the number  $n$ . Multiple choice



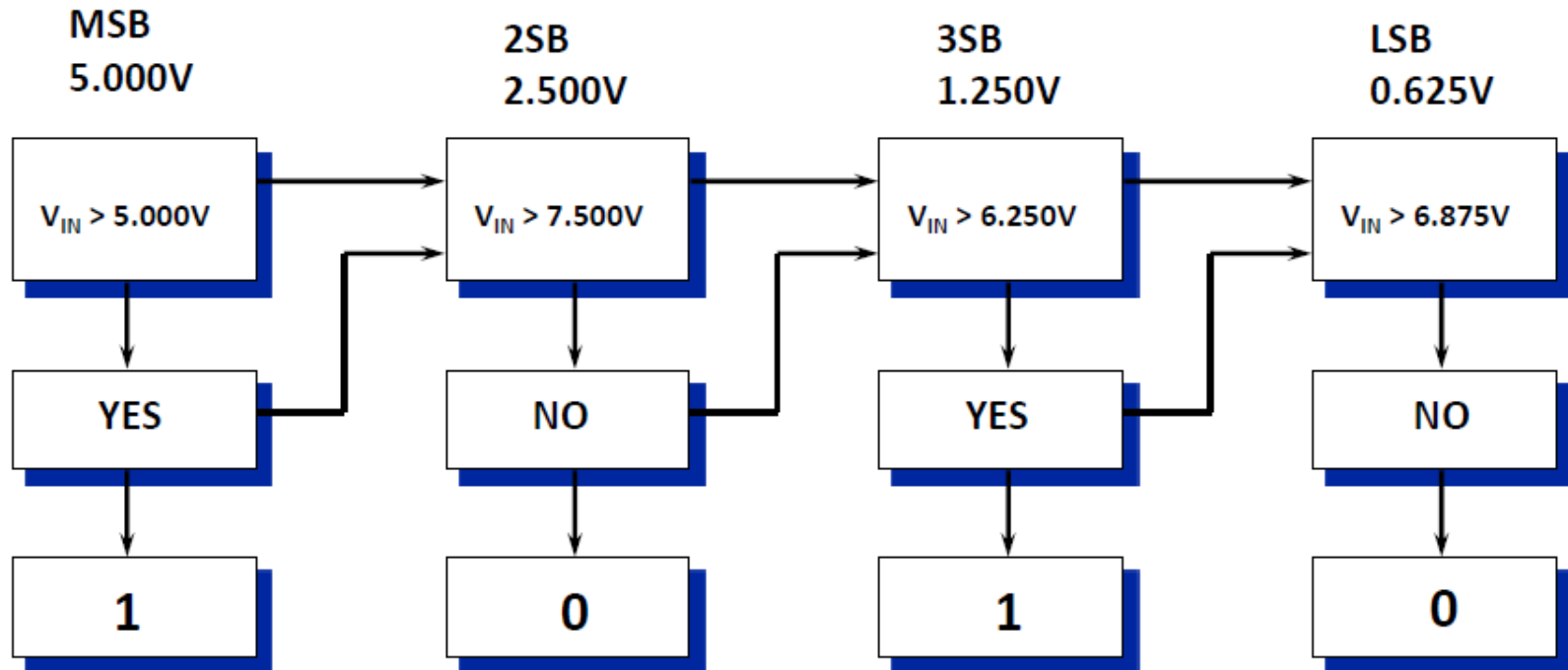
# Review of Electronics (Cont...)

400	200	100	0	0	0	0	0	0	0	No- Current
400	200	100	1	0	0	400	0	0	400	0.02
400	200	100	0	1	0	0	200	0	200	0.04
400	200	100	1	1	0	400	200	0	133.3333	0.06
400	200	100	0	0	1	0	0	100	100	0.08
400	200	100	1	0	1	400	0	100	80	0.1
400	200	100	0	1	1	0	200	100	66.66667	0.12
400	200	100	1	1	1	400	200	100	57.14286	0.14



# Review of Electronics (Cont...)

EXAMPLE : ANALOG INPUT = 6.428V, REFERENCE = 10.000V



# Assignment no. 1

In the design phase of the production line of fans. it required by the customer that fan to be controlled by three speeds (low, medium and high). your team leader asks you to identify the interface between the microcontroller and the fan motor. support your proposal by the circuit design.

Hint: motor driver, resistance ladder.

# Embedded Systems

- ~~Introduction to Embedded System.~~
- ~~Review of Electronics Part (KCL, KVL, Parallel/Series Resistance, DAC and ADC).~~
- Review for Digital Systems ( Binary Number, Logic, MOS Implantation, Computer Architecture).
- Intro to Programming Concepts (Structure and Concurrent).

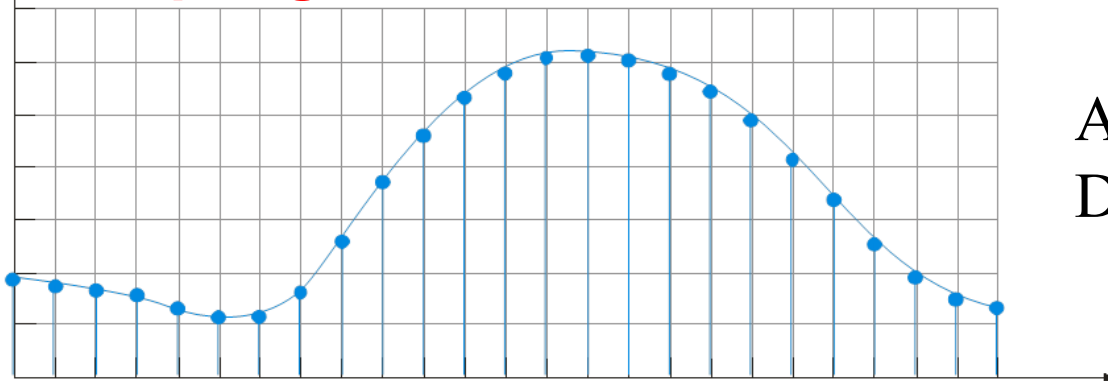


# Digital System

## Analog vs. Digital

Most natural quantities (such as temperature, pressure, light intensity, ...) are **analog** quantities that vary continuously.

### Sampling and Quantization



Analog = continuous  
Digital = discrete

**Digital systems** can process, store, and transmit data more efficiently but can only assign discrete values to each point.

# Binary System

- **Two discrete values:**
  - 1's and 0's
  - 1, TRUE, HIGH
  - 0, FALSE, LOW
- **1 and 0:** voltage levels, rotating gears, fluid levels, etc.
- Digital circuits use **voltage** levels to represent 1 and 0
- **Bit:** Binary digit
- **N-bit** binary number
  - How many values?  $2^N$  - Range:  $[0, 2^N - 1]$
  - Example: 3-digit binary number:
    - $2^3 = 8$  possible values - Range:  $[0, 7] = [000_2 \text{ to } 111_2]$

1's column  
2's column  
4's column  
8's column

$$1101_2 = 1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 13_{10}$$

one eight      one four      no two      one one

# Hexadecimal Numbers

- Base 16
- Shorthand for binary

Hex Digit	Decimal Equivalent	Binary Equivalent
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

# Bits, Bytes, Nibbles...

byte

10010110

nibble

10010110

most  
significant  
bit

least  
significant  
bit

CEBF9AD7

most  
significant  
byte

least  
significant  
byte

Binary

0011011011001101

Nibbles

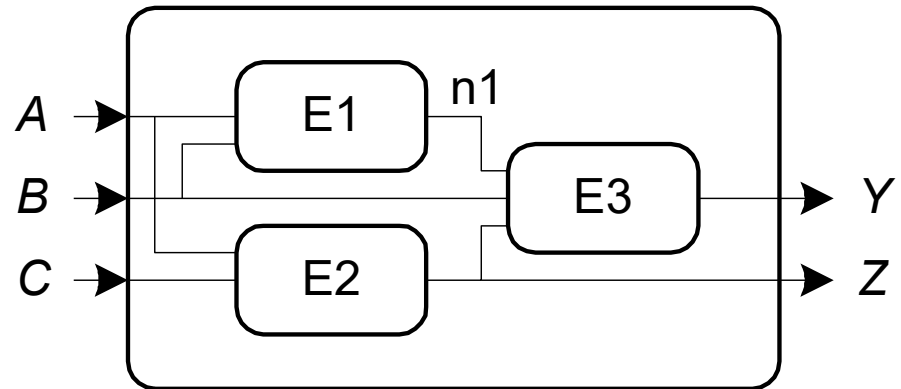
0011 0110 1100 1101

Hexadecimal

0x36CD

# Digital Circuit

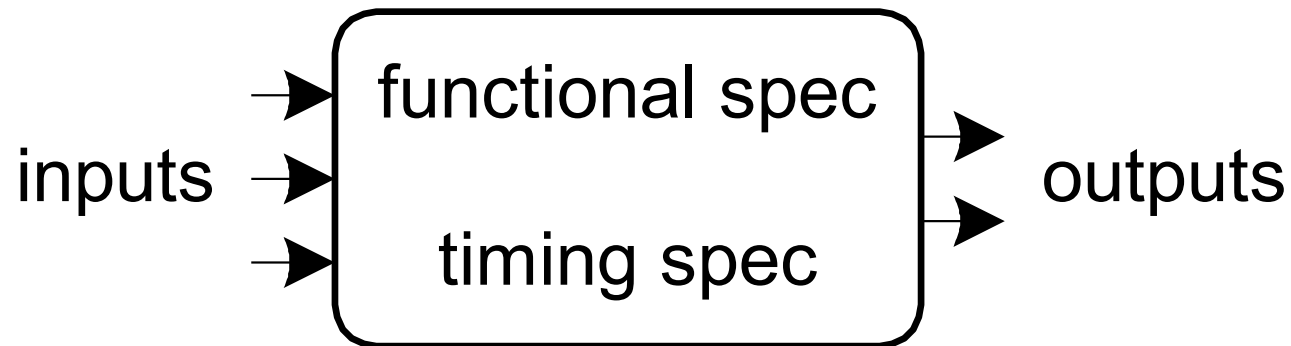
- Nodes
  - Inputs:  $A, B, C$
  - Outputs:  $Y, Z$
  - Internal:  $n1$
- Circuit elements
  - $E1, E2, E3$
  - Each a circuit



# Digital Circuit Specification

A logic circuit is composed of:

- Inputs
- Outputs
- Functional specification
- Timing specification



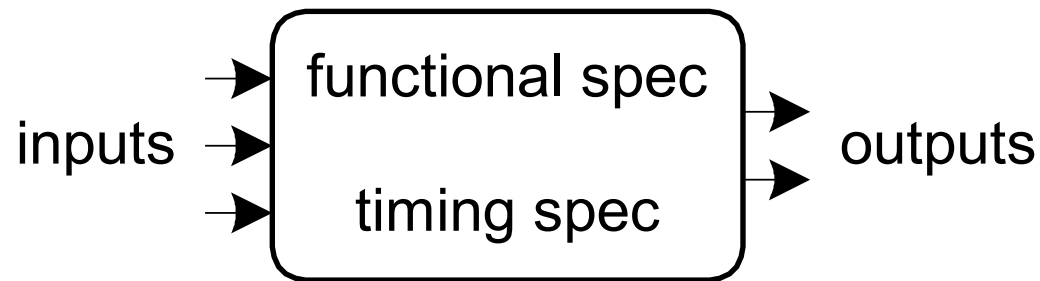
# Digital Circuit Types

- **Combinational Logic**

- Memoryless
- Outputs determined by current values of inputs

- **Sequential Logic**

- Has memory
- Outputs determined by previous and current values of inputs



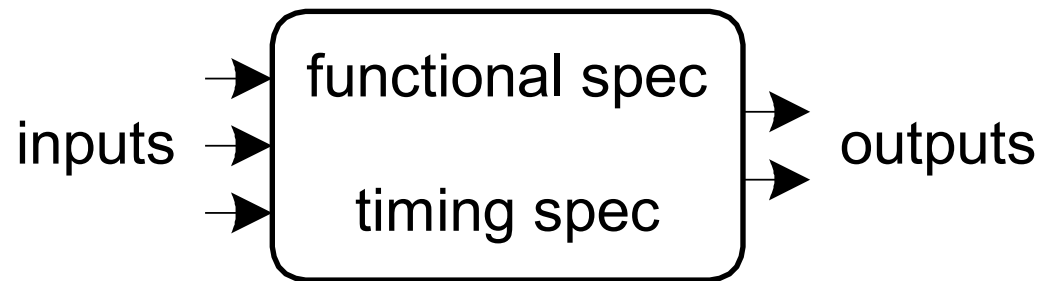
# Digital Circuit Types

- **Combinational Logic**

- Memoryless
- Outputs determined by current values of inputs

- **Sequential Logic**

- Has memory
- Outputs determined by previous and current values of inputs





# Definitions

- Complement: variable with a bar over it  
 $\bar{A}, \bar{B}, \bar{C}$
- Literal: variable or its complement  
 $A, \bar{A}, B, \bar{B}, C, \bar{C}$
- Implicant: product of literals  
 $\bar{A}\bar{B}\bar{C}, AC, BC$
- Minterm: **product** that includes all input variables  
 $\bar{A}\bar{B}\bar{C}, \bar{A}\bar{B}C, \bar{A}B\bar{C}, \bar{A}BC, A\bar{B}\bar{C}, A\bar{B}C, AB\bar{C}, ABC$
- Maxterm: **sum** that includes all input variables  
 $(A+\bar{B}+C), (\bar{A}+B+\bar{C}), (\bar{A}+\bar{B}+C)$

# Sum-of-Products (SOP) Form

- All equations can be written in SOP form
- Each row has a **minterm**
- A minterm is a product (AND) of literals
- Each minterm is **TRUE for that row** (and only that row)
- Form function by **ORing minterms where the output is TRUE**
- Thus, a sum (OR) of products (AND terms)

<b>A</b>	<b>B</b>	<b>Y</b>	<b>minterm</b>	<b>minterm name</b>
0	0	0	$\bar{A} \bar{B}$	$m_0$
0	1	1	$\bar{A} B$	$m_1$
1	0	0	$A \bar{B}$	$m_2$
1	1	1	$A B$	$m_3$

$$Y = F(A, B) = \bar{A}B + AB = \Sigma(1, 3)$$

# Product-of-Sums (POS) Form

- All Boolean equations can be written in POS form
- Each row has a **maxterm**
- A maxterm is a sum (OR) of literals
- Each maxterm is **FALSE for that row** (and only that row)
- Form function by **ANDing the maxterms for which the output is FALSE**
- Thus, a product (AND) of sums (OR terms)

<b>A</b>	<b>B</b>	<b>Y</b>	<b>maxterm</b>	<b>maxterm name</b>
0	0	0	$A + B$	$M_0$
0	1	1	$A + \overline{B}$	$M_1$
1	0	0	$\overline{A} + B$	$M_2$
1	1	1	$\overline{A} + \overline{B}$	$M_3$

$$Y = F(A, B) = (A + B)(\overline{A} + B) = \Pi(0, 2)$$

# Boolean Equations Example

- You are going to the cafeteria for lunch
  - You won't eat lunch ( $\bar{E}$ )
  - If it's not **open** ( $\bar{O}$ ) or
  - If they **only serve chicken** ( $C$ )
- Write a truth table for determining if you will eat lunch ( $E$ ).

$O$	$C$	$E$
0	0	0
0	1	0
1	0	1
1	1	0

# SOP & POS Form

- SOP – **sum-of-products**

<i>O</i>	<i>C</i>	<i>E</i>	minterm
0	0	0	$\bar{O} \bar{C}$
0	1	0	$\bar{O} C$
1	0	1	$O \bar{C}$
1	1	0	$O C$

$$E = O\bar{C}$$

$$= \Sigma(2)$$

- POS – **product-of-sums**

<i>O</i>	<i>C</i>	<i>E</i>	maxterm
0	0	0	$O + C$
0	1	0	$O + \bar{C}$
1	0	1	$\bar{O} + C$
1	1	0	$\bar{O} + \bar{C}$

$$E = (O + C)(O + \bar{C})(\bar{O} + \bar{C})$$

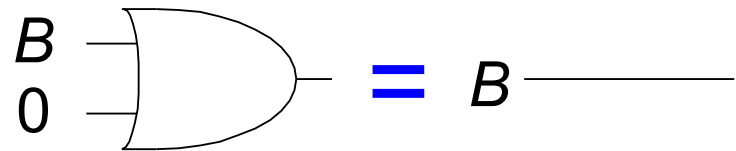
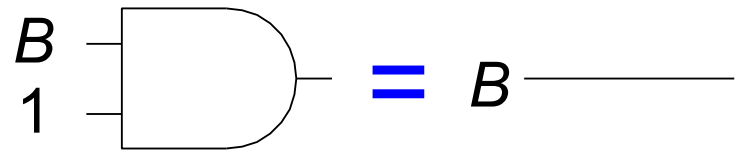
$$= \Pi(0, 1, 3)$$

# Boolean Algebra

- Axioms and theorems to **simplify** Boolean equations
- Like regular algebra, but simpler: variables have only two values (1 or 0)
- **Duality** in axioms and theorems:
  - ANDs and ORs, 0's and 1's interchanged

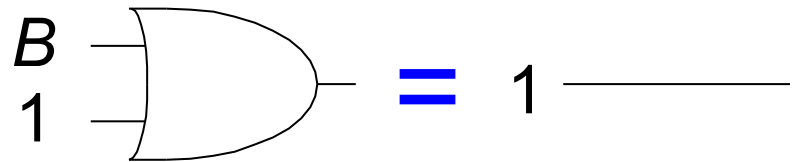
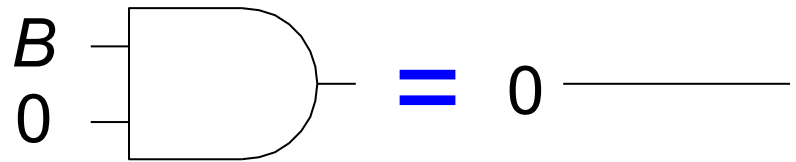
# T1: Identity Theorem

- $B \cdot 1 = B$
- $B + 0 = B$



# T2: Null Element Theorem

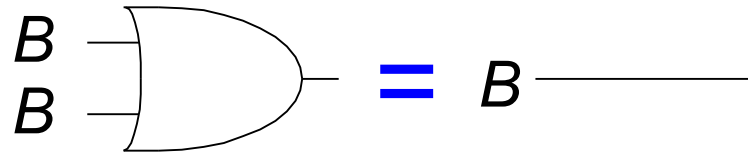
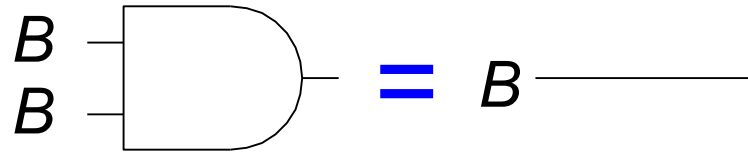
- $B \cdot 0 = 0$
- $B + 1 = 1$





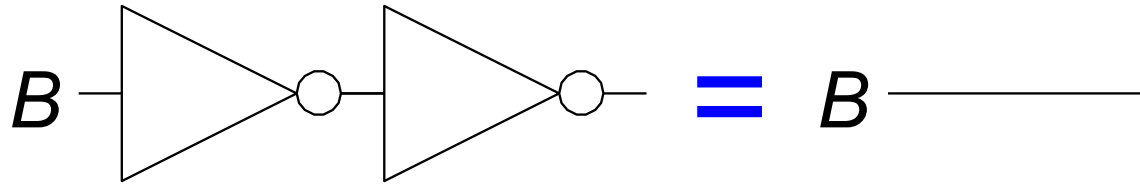
# T3: Idempotency Theorem

- $B \cdot B = B$
- $B + B = B$



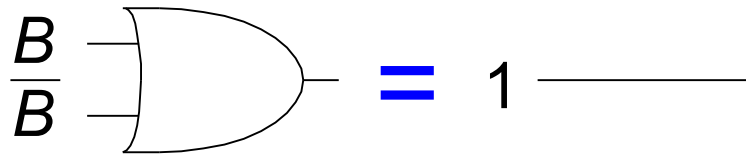
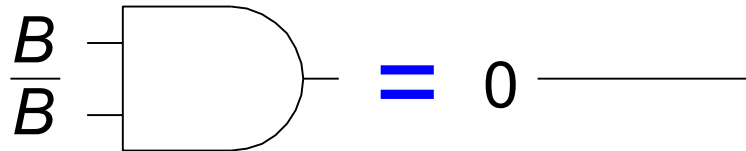
# T4: Identity Theorem

- $\overline{\overline{B}} = B$



# T5: Complement Theorem

- $B \cdot \bar{B} = 0$
- $B + \bar{B} = 1$



# Boolean Theorems Summary

	Theorem		Dual	Name
T1	$B \bullet 1 = B$	T1'	$B + 0 = B$	Identity
T2	$B \bullet 0 = 0$	T2'	$B + 1 = 1$	Null Element
T3	$B \bullet B = B$	T3'	$B + B = B$	Idempotency
T4		$\overline{\overline{B}} = B$		Involution
T5	$B \bullet \overline{B} = 0$	T5'	$B + \overline{B} = 1$	Complements

# Boolean Theorems of Several Vars

Theorem		Dual		Name
T6	$B \bullet C = C \bullet B$	T6'	$B + C = C + B$	Commutativity
T7	$(B \bullet C) \bullet D = B \bullet (C \bullet D)$	T7'	$(B + C) + D = B + (C + D)$	Associativity
T8	$(B \bullet C) + B \bullet D = B \bullet (C + D)$	T8'	$(B + C) \bullet (B + D) = B + (C \bullet D)$	Distributivity
T9	$B \bullet (B + C) = B$	T9'	$B + (B \bullet C) = B$	Covering
T10	$(B \bullet C) + (B \bullet \overline{C}) = B$	T10'	$(B + C) \bullet (B + \overline{C}) = B$	Combining
T11	$(B \bullet C) + (\overline{B} \bullet D) + (C \bullet D)$ $= B \bullet C + \overline{B} \bullet D$	T11'	$(B + C) \bullet (\overline{B} + D) \bullet (C + D)$ $= (B + C) \bullet (\overline{B} + D)$	Consensus
T12	$\overline{B_0 \bullet B_1 \bullet B_2 \dots}$ $= (\overline{B_0} + \overline{B_1} + \overline{B_2} \dots)$	T12'	$\overline{B_0 + B_1 + B_2 \dots}$ $= (\overline{B_0} \bullet \overline{B_1} \bullet \overline{B_2})$	De Morgan's Theorem

**Note:** T8' differs from traditional algebra: OR (+) distributes over AND (•)

# Simplifying Boolean Equations

**Example 1:**

$$Y = AB + \bar{A}B$$

$$= B(A + \bar{A}) \quad \text{T8}$$

$$= B(1) \quad \text{T5'}$$

$$= B \quad \text{T1}$$

# Simplifying Boolean Equations

**Example 2:**

$$Y = A(AB + ABC)$$

$$= A(AB(1 + C)) \quad \text{T8}$$

$$= A(AB(1)) \quad \text{T2'}$$

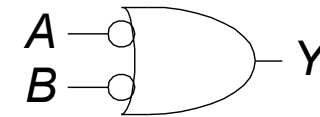
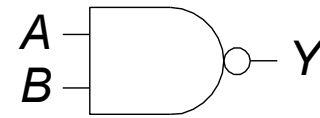
$$= A(AB) \quad \text{T1}$$

$$= (AA)B \quad \text{T7}$$

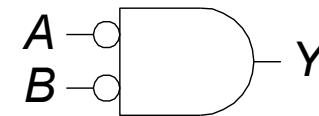
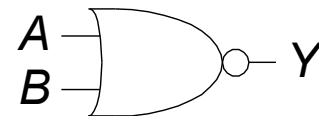
$$= AB \quad \text{T3}$$

# DeMorgan's Theorem

- $Y = \overline{AB} = \overline{A} + \overline{B}$



- $Y = \overline{A + B} = \overline{A} \cdot \overline{B}$

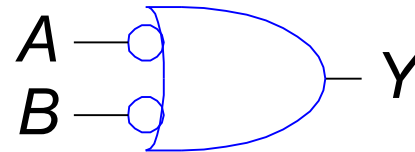
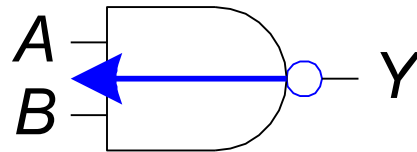




# Bubble Pushing

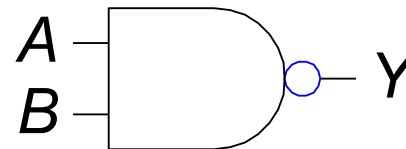
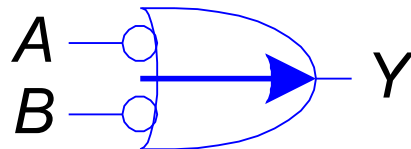
- **Backward:**

- Body changes
- Adds bubbles to inputs



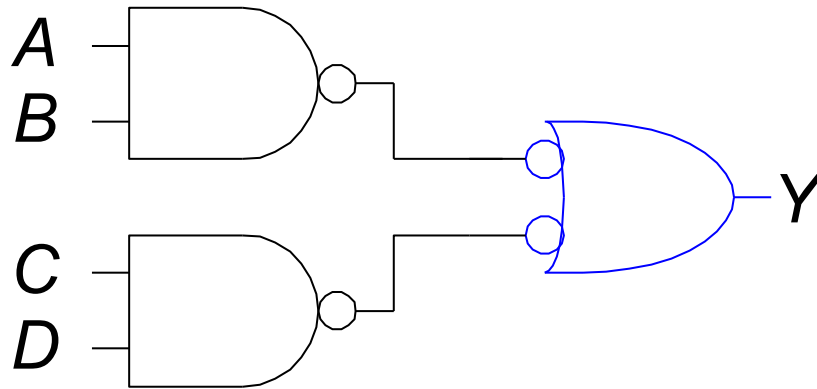
- **Forward:**

- Body changes
- Adds bubble to output



# Bubble Pushing

- What is the Boolean expression for this circuit?



$$Y = AB + CD$$

# Boolean Algebra

- $A \& B = B \& A$  ... Commutative Law
- $A | B = B | A$  ... Commutative Law
- $(A \& B) \& C = A \& (B \& C)$  ... Associative Law
- $(A | B) | C = A | (B | C)$  ... Associative Law
- $(A | B) \& C = (A \& C) | (B \& C)$  ... Distributive Law
- $(A \& B) | C = (A | C) \& (B | C)$  ... Distributive Law
- $\sim(A | B) = (\sim A) \& (\sim B)$  ... De Morgan's Theorem
- $\sim(A \& B) = (\sim A) | (\sim B)$  ... De Morgan's Theorem

$A \& 0 = 0$  ... Identity of 0

$A | 0 = A$  ... Identity of 0

$A \& 1 = A$  ... Identity of 1

$A | 1 = 1$  ... Identity of 1

$A | A = A$  ... Property of OR

$A | (\sim A) = 1$  ... Property of OR

$A \& A = A$  ... Property of AND

$A \& (\sim A) = 0$  ... Property of AND

$\sim(\sim A) = A$  ... Inverse

# Operators (Bitwise / logical)

A	B	C	A & (B   C)	A && (B    C)
01	01	11	01	True
00	01	11	00	False

# Signed Binary Numbers

- Sign/Magnitude Numbers
  - 1 sign bit,  $N-1$  magnitude bits
  - Positive number: sign bit = 0 , Negative number: sign bit = 1
  - Example, 4-bit sign/mag representations of  $\pm 6$ :
    - +6 = **0110**
    - 6 = **1110**
  - Range of an N-bit sign/magnitude number:  **$[-(2^{N-1}-1), 2^{N-1}-1]$**
- Two's Complement Numbers
  - The most significant bit still indicates the sign (1 = negative, 0 = positive)
  - Range of an N-bit two's comp number:  **$[-(2^{N-1}), 2^{N-1}-1]$**

1. **1001**

2. **+ 1**

**$1010_2 = -6_{10}$**

$$\begin{array}{r} 111 \\ 0110 \\ + 1010 \\ \hline 10000 \end{array}$$

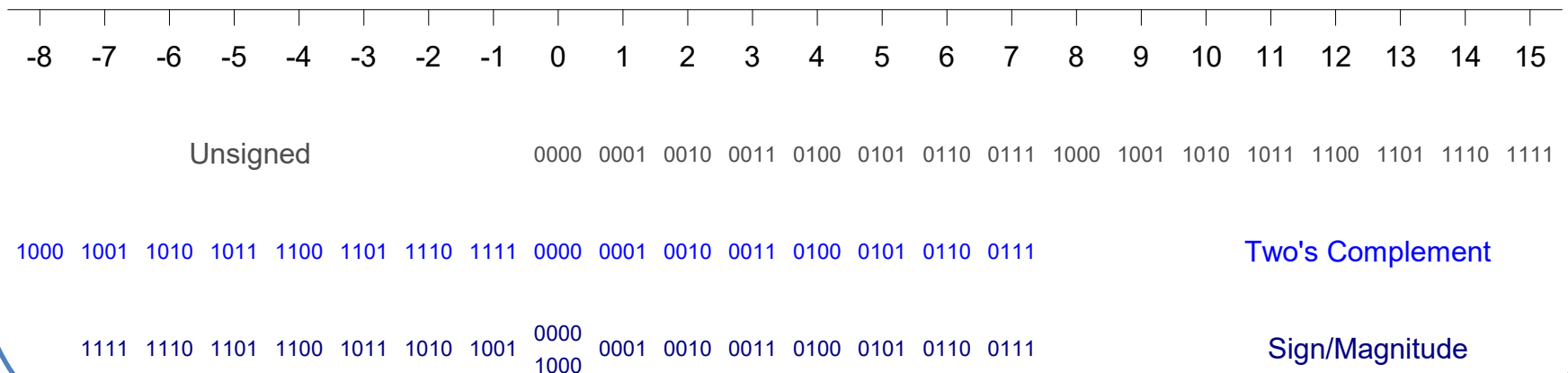
# Signed Binary Numbers (Cont...)

- Sign Extension
  - Sign bit copied to msb's
  - Number value is same
- **Example 1:**
  - 4-bit representation of 3 = 0011
  - 8-bit sign-extended value: 00000011
- **Example 2:**
  - 4-bit representation of -5 = 1011
  - 8-bit sign-extended value: 11111011

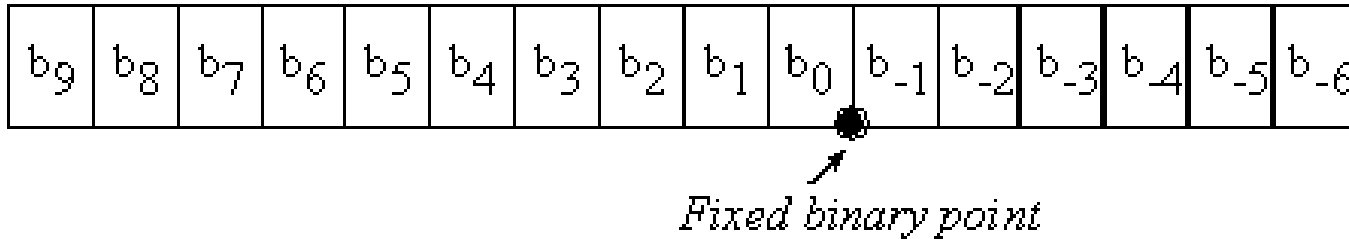
# Signed Binary Numbers (Cont...)

Number System	Range
Unsigned	$[0, 2^N-1]$
Sign/Magnitude	$[-(2^{N-1}-1), 2^{N-1}-1]$
Two's Complement	$[-2^{N-1}, 2^{N-1}-1]$

For example, 4-bit representation:



# Fixed Point Binary Numbers



- **fixed-point numbers:** express values to the computer **non-integer** values.
- A **fixed-point number** contains two parts: variable **integer**, called  $I$ . fixed **constant**, called the **resolution**.
- The fixed constant will **NOT be stored on the computer**. The fixed constant is something we keep track of while designing the software operations. !!!
- The **precision** of a number system is the **total number of distinguishable values** that can be represented.
- The precision of a fixed-point number is **the number of bits used to store the variable integer**.

$$\text{Binary fixed-point value} = I \cdot 2^n$$



# Fixed Point Binary Numbers

01101100

0110.1100

$$2^2 + 2^1 + 2^{-1} + 2^{-2} = 6.75$$

# Addition

- Decimal

$$\begin{array}{r} 11 \leftarrow \text{carries} \\ 3734 \\ + 5168 \\ \hline 8902 \end{array}$$

- Binary

$$\begin{array}{r} 11 \leftarrow \text{carries} \\ 1011 \\ + 0011 \\ \hline 1110 \end{array}$$

- Digital systems operate on a **fixed number of bits**
- **Overflow**: when result is too big to fit in the available number of bits

$$\begin{array}{r} 111 \\ 1011 \\ + 0110 \\ \hline 10001 \end{array}$$

Overflow!

[Ariane-5 Rocket Explosion \(2002\)](#)

# Multiplication

## Decimal

$$\begin{array}{r} 230 \\ \times 42 \\ \hline 460 \\ + 920 \\ \hline 9660 \end{array}$$

$$230 \times 42 = 9660$$

multiplicand

multiplier

partial  
products

result

## Binary

$$\begin{array}{r} 0101 \\ \times 0111 \\ \hline 0101 \\ 0101 \\ 0101 \\ + 0000 \\ \hline 0100011 \end{array}$$

$$5 \times 7 = 35$$

- Partial products formed by **multiplying a single digit** of the multiplier with multiplicand
- **Shifted partial products** summed to form result